#### EBMT Based on Finite Automata State Transfer Generation

# Feiliang Ren renfeiliang@gmail.com



## Contents

> Introduction

Related researches

System Structure of Our CJ EBMT System

- Generation Based on Finite Automata State Transfer
  - Building Links

State Assignments

Translation Generation

> Experiments

Conclusions and Future Work

### Introduction

- EBMT : a method of translation by the principle of analogy
- > Three basic modules
  - Matching module
  - Alignment module
  - Recombination module
- > The last two modules can be regarded as a translation generation module.
  - Semantic-based generation approach
    - > Obtains an appropriate translation fragment for each part of the input sentence.
    - > Final translation is generated by recombining the translation fragments in some order.
    - > Shortcoming: doesn't take into account the fluency between the translation fragments
  - Statistical approach
    - > Selects translation fragments with a statistical model
    - Can improve the fluency between the translation fragments by using n-gram co-occurrence statistics.
    - Shortcoming: doesn't take into account the semantic relation between the example and the input sentence
  - > Method based on tree string correspondence (TSC) and statistical generation
    - Can solve the shortcomings of the above generation approaches ;
    - But: depends on the tree parser so much that if the parser doesn't work well, it is impossible to generate a proper translation result.

# System Structure of Our CJ EBMT System

- Our generation method
  - Uses the target sentence of the selected example to generate the translation of the input sentence.
  - Generate the translation in a finite automata state transfer manner.



Structure of our EBMT System

# Generation Based on Finite Automata State Transfer

> Matching: select translation examples for the input sentence

Method: a combined method based on substantive word matching and stop word matching

WordSim(A,B)= $2 \cdot \frac{SameWC(A,B)}{len(A) + len(B)}$ 

StopWord\_Sim(A,B)=exp(( $abs(StopWord(A) - StopWord(B)) \times \beta$ )

 $final \_Sim(A,B) = \frac{WordSim(A,B)}{StopWord \_Sim(A,B)}$ 

Generation

- Step 1 、 Build links from the fragments in the input sentence to the fragments in the target sentence of the selected example
- Step 2 , Assign states to each of these links
- Step 3 Construct a finite automaton and generate the translation result in an automaton state transfer manner

# Step 1 for Generation: Building Links

- Link : a link from a fragment in one sentence  $S_1$  to a fragment in another sentence  $S_2$  is defined as a 3-tuple ( $Sf_i$ ,  $Tf_j$ , t).
  - $\succ$  Sf<sub>i</sub>: a fragment in S<sub>1</sub>
  - $\succ$   $Tf_i$ : a fragment in  $S_2$
  - t: link type, we define four link types: I, R, D, N, which mean inserting, replacing, deleting and outputting directly respectively
- > Build links from the fragments in the input sentence S to the fragments in the target sentence B of the selected example (A, B)
  - First: Build links from S's fragments to A's fragments using a revised edit distance

algorithm (will be shown in the next slide). Its result is denoted as  $LinkSet(S \rightarrow A)$ .

- Second: Build links from S's fragments to B's fragments (denoted as  $LinkSet(S \rightarrow B)$ ) according to following rules.
  - (a) For a link in *LinkSet(S→A)*, if neither its source fragment nor its target fragment is null, replace its target fragment with this target fragment's corresponding aligned fragment in *B*, and add this new link to *LinkSet(S→B)*.
  - (b) For a link in *LinkSet(S→A)* whose target fragment is null, add it to *LinkSet(S→B)* directly.
  - (c) For those fragments in B that have not been linked, build links for each of them by assigning a null source fragment and a D link type to them respectively, and add these links to LinkSet(S→B).
  - (d) Reorder the items of  $LinkSet(S \rightarrow B)$  in their target fragments' order in sentence B Feiling Ren 2007-10

# Step 1 for Generation: Building Links

The algorithm for building links from S's fragments to A's fragments is shown

```
as followings.
m = length(S_1), n = length(S_2)
d[0][0] = 0; tags[0][0] = 0;
for i=1 to m
   d[i][0]=q+d[i-1][0]; tags[i][0]='D'
for j=1 to n
   d[0][i]=r+d[0][i-1]; tags[0][i]='I'
for i=1 to m
  for j=1 to n
    p = computeCost(S_{1}[i-1], S_{2}[j-1]);
    a = d[i-1][j-1] + p;
    b=d[i-1][j] + q;
    c=d[i][j-1] + r;
    d[i][j] = min(a,b,c);
    if(min=a and p==0)
        tags[i][j] = N';
    else if (min=a)
        tags[i][j] = R';
     else if (min=b)
   Revised Edit Distance Algorithm
```

*computeCost* is a function to compute two fragments' linking cost based on their lexical forms and their head words' POSs.

- If two fragments' lexical forms are the same and their head words' POSs are the same too, this cost is zero;
- if two fragments' lexical forms are the same but their head words' POSs are different, this cost is 0.2;
- otherwise, this value is assigned by human's experiences according to the two fragments' head words' POSs as shown in the following table

$PosPair(c_i, c_j)$	W <sub>i</sub>
(noun, noun)	0.5
(noun, auxiliary)	0.8
(noun, adjective)	0.85

#### Linking Cost for Two Fragments

### Step 1 for Generation: Building Links

> The whole process of this step can be shown in the following figure



### Step 2 for generation: States Assignment

- States for Non-I Type's Links
  - > If its link type is R, a state named  $S_R$  is assigned
  - $\succ$  If its link type is *D*, a state named *S\_D* is assigned;
  - $\succ$  If its link type is *N*, a state named *S*\_*N* is assigned.
- > States for *I* Type's Links
  - Consider context of current *I*-type link's pre- and post- links
  - Consider link shapes
  - Define 12 basic link shapes and 3 extended link shapes for *I*-type link, and map each of these link shapes to an *I*-type link's state.

### Step 2 for generation: States Assignment

➢ Basic States for *I*-type's Link



# Step 2 for generation: States Assignment



#### > Extended states can be converted into basic states

- For state 13, move rightward until find a non-*I* type's link, if this link's target fragment is null, convert it to state 6; otherwise, convert it to a state among state 1 to state 5 according to the link shapes of fragment *i*-1's link and the new found link; if can't find a non-*I* type's link in current link's right side, convert it to state 11.
- For state 14, move rightward until find a non-*I* type's link, if this link's target fragment is null, convert it to state 8, otherwise, convert it to state 7; if can't find a non-*I* type's link in current link's right side, convert it to state 12.
- For state 15, move rightward until find a non-*I* type's link, if this link's target fragment is null, convert it to state 10, otherwise, convert it to state 9; if can't find a non-*I* type's link in current link's right side, move leftward until find a non-*I* type's link (this link will be found always) and convert it to state 11.

### Step 3 for generation: Translation Generation

- Generation Operation for Non-I Type Links' States
  - If a link's state is S\_R, replace this link's target fragment with its source fragment's translation, and denote this operation as O(R);
  - If a link's state is S\_D, delete this link's target fragment, and denotes this operation as O(D);
  - ➢ If a link's state is S\_N, remain this link's target fragment unchanged, and denote this operation as O(N).
- Generation Operation for I Type Links' States
  - Take its source fragment's pre- and post- fragments into account and judge: whether the fragment combinations (*i*-1,*i*,*i*+1), (*i*-1,*i*) and (*i*,*i*+1) are chunks. If they are chunks, look up their corresponding translations in dictionary, otherwise, look up *i*'s translation in dictionary (we assume its translation can be found always).

According to current *I*-type link's state and the recognized chunk information, we choose one of these chunks as current *I*-type link's new source fragment for later processing, and define 10 possible generation

2007 - 10

**Feiliang Ren** 

# Step 3 for generation: Translation Generation

#### ➢ Generation Operation for *I* Type Links' States

- O(0): Delete the links that take *B*'s fragments among m+1 to *n* as their target fragments. And for the link that takes *B*'s fragment *m* as target fragment, replace *m* with the translation of current *I*-type link's new source fragment.
- **O(1)**: For the link that takes *B*'s fragment *m* as target fragment, replace *m* with the translation of current *I*-type link's new source fragment.
- *O(2):* For the link that takes *B*'s fragment *n* as target fragment, replace *n* with the translation of current *I*-type link's new source fragment.
- *O(3):* For the link that takes *B*'s fragment *m* as target fragment, add the translation of current *I*-type link's new source fragment to the end of *m*.
- *O(4):* For the link that takes *B*'s fragment *n* as target fragment, add the translation of current *I*-type link's new source fragment to the end of *n*.
- **O(5):** For the link that takes *B*'s fragment *m* as target fragment, replace *m* with the translation of current *I*-type link's new source fragment. And delete the link that takes *B*'s fragment *n* as target fragment.
- **O(6):** For the link that takes *B*'s fragment *n* as target fragment, replace *n* with the translation of current *I*-type link's new source fragment. And delete the link that takes *B*'s fragment *m* as target fragment.
- *O(7)*: For the link that takes *B*'s fragment *m* as target fragment, add the translation of current *I*-type link's new source fragment before *m*.
- **O(8):** For the link that takes *B*'s fragment *n* as target fragment, add the translation of current *I*-type link's new source fragment before *n*.
- **O(9):** Do not modify any link's target fragment.

# Step 3 for generation: Translation Generation

▷ Based on  $LinkSet(S \rightarrow B)$  and the assigned states, we construct an automaton that has a similar form as shown in the following figure

![](_page_13_Figure_4.jpeg)

- $\succ$  B is a start state
- $\succ$  E is an end state
- $\succ$  {*I*, *R*, *D*, *N*} are link types
- >  $\{O(N), O(D), O(R)\}$  in parallelogram are the operations
- # is a fictitious symbol that indicates the end of the automaton's input
- {S\_R, S\_D, S\_N} are states correspond to non-I type's links
- S\_I' is a state set that corresponds to I-type's links

#### Step 3 for generation: Translation Generation

![](_page_14_Figure_3.jpeg)

- O' in the operation of state 3 means the automaton generates the fragment combination (*i*-1,*i*,*i*+1)'s translation by simply joining their single fragment's translations together.
- d<sub>1</sub> means the semantic distance from fragment *i* to fragment
   *i*-1, and d<sub>2</sub> means the semantic distance from fragment *i* to fragment *i*+1, and they are computed as following formula:

$$dist(f_1, f_2) = \sum_{c_i \in f_1} \sum_{c_j \in f_2} w_k(PosPair(c_i, c_j))$$

![](_page_14_Figure_7.jpeg)

**Feiliang Ren** 

2007 - 10

#### Step 3 for generation: An Example

- ▶ Suppose S is "他很爱他的妻子(He loves his wife very much)". The selected example (A,B) is "(他爱他的妈妈(He loves his mother),彼は、彼の母を愛し ています(He loves his mother))".
- ▶ After building links, LinkSet(S→B) is: (他(he), 彼(he), N), (null, は(ha), D), (很(very much), null, I), (他的(his), 彼の(his), N), (妻子(wife), 母(mother), R), (null, を(wo), D) (爱(loves), 愛しています(loves), N)
- ▶ Its corresponding state sequence is: S\_N, S\_D, S\_I\_4(the forth state in the basic *I*-type's links), S\_N, S\_R, S\_D, S\_N.
- Construct an automaton, and begin to states transfer and t ranslation generation.
- For the link ( $(\underline{h}(he), \underline{h}(he), N)$ ), its state is  $S_N$ . The automaton executes operation O(N) and does not mo dify this link's target fragment.
- For the link (null, (ha), D), its state is S\_D. The automaton executes operation O(D) and deletes this link's target fragment.
- ▶ For the link (很(very much), null, I), its state is S\_I\_4. If the fragment combination (i-1,i) "他 很(he...very much)" is a chunk and the corresponding translation
- is "彼は、とても(he...very much)", the automaton executes operation O(1). It first takes this recognized chunk as current link's new source fragment. Then it
- selects the link whose target fragment is "彼(he)", and this link is (他(he),彼(he),N). Thirdly, it replaces the selected link's target fragment with the
- translation of current *I*-type link's new source fragment. At last the selected link is changed to (他(he), 彼は、とても(he…very much), N).
- ト For the link (他的(*his*),彼の母(*his*),N), its state is  $S_N$ . The automaton executes oper ation O(N) and does not modify this link's target fragment.
- ▶ For the link (妻子(*wife*),母(*mother*),R), its state is S\_R. The automaton executes operation O(R) and replaces this link's target fragment with its source fragment's translation. Finally current link is changed to (妻子(*wife*),妻(*wife*),R).
- For the link (null,  $\mathcal{E}(wo), D$ ), its state is S\_D. The automaton executes operation O(D) and deletes this link's target fragment.
- ▶ For the link (爱(loves),愛しています(loves),N), its state is S\_N. The automaton executes operation O(N) and does not modify this link's target fragment.
- ▶ At last, the automaton ends the state tran sfer process and outputs *LinkSet(S→B)*'s modified target fragment s equence "彼は、とても彼の妻愛しています(he loves his wife very much)" and takes it as the input sentence 's translation.

#### Experiments

#### System Resources

- Bilingual Corpus: We collect 10083 Chinese-Japanese bilingual sentences from Internet in Olympic domain as examples
- Bilingual Dictionary: A bilingual dictionary is used to translate the input fragment and to judge whether an input fragment is a chunk.
- Language Model: We collected an approximate 1,400,000 words' Japanese monolingual corpus and a similar size's Chinese monolingual corpus from Internet, and trained a standard trigram Japanese language model for Chinese-to-Japanese EBMT system and a standard trigram Chinese language model for Japanese-to-Chinese EBMT system respectively.
- Test Corpus: We collect another 100 bilingual sentences in Olympic domain from Internet as test corpus.

#### Experimental Result

Experimental Results for Chinese-to-Japanese EBMT

Suctam

Method	NIST	BLEU
Baseline	4.8321	0.4913
Our System	5.9729	0.7705

Experimental Results for Japanese-to-Chinese EBMT

Su	stem	
U y	Stom	

Method	NIST	BLEU
Baseline	4.1275	0.4076
Our System	5.0976	0.5908

#### Experiments----Some Translation Examples

Input:	我们的足球被对方前锋拦截
Output:	私 たち の サッカー は 相手 の 前鋒 に 阻まれ た
Input:	摔跤强国俄罗斯和日本有很多足球俱乐部
Output:	レスリング の 強国 ロシア と 日本 に は 很多 足球 俱乐部 が ある
Input:	中国运动员孙英杰今年一直主攻马拉松
Output:	中国 の スポーツ 選手 英傑 は 今年 ずっと マラソン を 専攻 として いる
	Same Translation Develts for Chinese to Lengeners Translation
	Some Translation Results for Chinese-to-Japanese Translation
Input:	客判員はいささかのためらいもなくペナルティーキックを科した
Input: Output:	審判員はいささかのためらいもなくペナルティーキックを科した 裁判 毫不犹豫 地 判罚 点球
Input: Output: Input:	審判員はいささかのためらいもなくペナルティーキックを科した 裁判 毫不犹豫 地 判罚 点球 中国のチームにはマンツーマンディフェンス戦術がある
Input: Output: Input: Output:	審判員はいささかのためらいもなくペナルティーキックを科した 裁判 毫不犹豫 地 判罚 点球 中国のチームにはマンツーマンディフェンス戦術がある 中国 的 队 有 町人 战术
Input: Output: Input: Output: Input:	Some Translation Results for Chinese-to-Japanese Translation        審判員はいささかのためらいもなくペナルティーキックを科した        裁判 毫不犹豫 地 判罚 点球        中国のチームにはマンツーマンディフェンス戦術がある        中国 的 队 有 盯人 战术        スウェ - デンのチ - ム20分のペナルティを受けた
Input: Output: Input: Output: Input: Output:	Some Translation Results for Chinese-to-Japanese Translation        審判員はいささかのためらいもなくペナルティーキックを科した        裁判 毫不犹豫 地 判罚 点球        中国のチームにはマンツーマンディフェンス戦術がある        中国 的 队 有 盯人 战术        スウェ - デンのチ - ム20分のペナルティを受けた        瑞典 队 被 罚 了 20 分

Some Translation Results for Japanese-to-Chinese Translation

# Conclusions and Future Work

#### Conclusions:

- > The natural of the states are some transfer rules.
- > Our work can work on most of language pairs.
- It doesn't need any complicated parsers.
- Future Work
  - Merge syntax analysis into our method
  - > Merge probability knowledge into state assignment and generation.

# The End

≻Thanks!

If you have any question, please contact me by renfeiliang@gmail.com, or renfeiliang@ise.neu.edu.cn

Welcome to my website: http://www.nlplab.cn/renfeiliang/

![](_page_19_Picture_6.jpeg)